

Abstraction of Concurrency and Communication in VHDL

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VHDL in the design flow

- At the system level, specify aspects of the system
 - structure (architectural partitioning)
 - behaviour (data transforms, reactions)
 - . . .
- VHDL allows description of structure and behaviour
- Prefer single language approach
 - system level to detailed design (hardware/software)
 - aids refinement
 - avoid semantic mismatch

Problem with VHDL

- Behaviour includes concurrency and communication
- Communication performed using signals
 - assignment, resolution, waits
- Signals do not express synchronization
 - *cf.* system description and programming languages
 - assignment schedules a message for a particular time
 - receiver may miss the message
 - if it does not respond at that time

Ad hoc solutions

- Communication using handshaking protocol
 - requires extra signals
 - requires protocol implementation
- Explicitly instantiated message queues
 - extra components
- Both detract from abstraction of communication
 - introduction of extraneous artefacts

Previous extension proposals

- Vista OO-VHDL (Swamy *et al*)
- LaMI (Benzakki & Djaffri)
- Bournemouth/IBM-UK (Cabanis *et al*)
- ➔ Based on OO extensions to VHDL
 - entity/architecture as a class, components as objects
 - entity specifies operations invoked by other processes
 - architecture encapsulates state
- ⇒ Monitor semantics with concurrency control

Other description languages

for describing behaviour, eg:

- StateCharts
- Estelle
- SDL
- CSP
- . . .
- ➔ Concurrent processes with message passing

StateCharts

- Extended FSM
 - hierarchical states
 - concurrent composition of substates
 - transitions on *events* (optionally *guarded*)
 - communication:
 - action in one chart \Rightarrow event in another chart
- Used in UML
 - events can be parameterized
 - procedural actions
- States are statically specified

Estelle

- Hierarchy of extended NFAs (*tasks*)
 - instances of declared modules
 - can be statically or dynamically instantiated
 - transitions enabled by event occurrence (msg arrival)
 - actions on transitions are Pascal-like statements
 - lock-step parallelism (select-transition/fire/action)
- Communication: message passing via links
 - links can be statically or dynamically created
 - non-blocking send, buffered at receiver

SDL

- Hierarchy of statically specified process sets
- Processes: statically or dynamically instantiated
 - instances of declared process types
 - behaviour specified as extended FSM
 - transitions enabled by message arrival
 - actions: assignment, branching, process creation, procedure call, message send
- Communication: message passing via channels
 - channels between process sets
 - non-blocking send, buffered at receiver
- Also a form of RPC

CSP

- Statically specified processes
 - contain sequentially executed actions
 - variable assignment, send, receive
- Synchronous message passing
 - statically specified channels

Comparison with VHDL

- Concurrent process model with message passing
 - ≡ concurrent FSMs communicating by events
- VHDL
 - statically specified processes & channels (signals)
 - asynchronous unbuffered message passing (!?)
- Alternatives
 - static vs. dynamic process instantiation
 - static vs. dynamic channel creation
 - sync vs. async message passing

Language design issues

- Message passing mechanism and semantics
- Communication abstraction in entity interface
- Dynamic process creation/termination semantics
 - Process abstraction and interface
- Integration with existing language
- Integration with other extension (eg, SUAVE)

Message passing mechanism (1)

- Alternative 1: introduce new mechanism
 - eg, message channel, send & receive operations

```
channel elevator_call : floor_number;  
channel elevator_location : floor_number;  
  
elevator : process is  
begin  
    ...  
    receive calling_floor from elevator_call;  
    send current_floor to elevator_location;  
    ...  
end process elevator;
```

Communication in entity interface (1)

- Alternative 1

```
entity operator_console is  
    port ( channel status : in status_msg;  
          channel command : out command_msg );  
end entity operator_console;
```

Message passing mechanism (2)

- Alternative 2: generalize existing mechanism
 - eg, abstract signals
 - transactions queued without delay mechanism
 - wait receives next value from queue

```
signal elevator_call : floor_number abstract;  
signal elevator_location : floor_number abstract;  
  
elevator : process is  
begin  
    ...  
    wait on elevator_call;  -- receive next message  
    calling_floor := elevator_call;  
    elevator_location <= current_floor;  -- send message  
    ...  
end process elevator;
```

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Communication in entity interface (2)

- Alternative 2

```
entity operator_console is  
    port ( signal status : in status_msg abstract;  
          signal command : out command_msg abstract );  
end entity operator_console;
```

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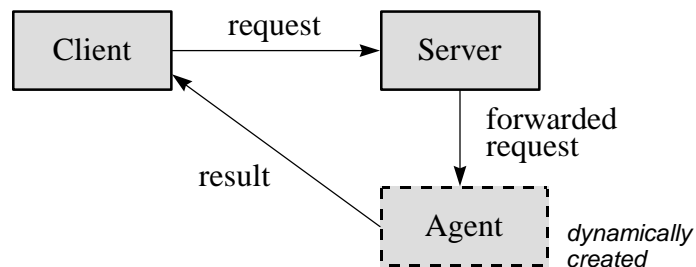
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Dynamic process creation

- Eg, for multithreaded server
- Process types
 - needed for dynamic instantiation
 - parameterized by channel/signal
- Channel/signal types
 - needed for dynamic creation of channel/signal
 - for communication with dynamically created process

Process & channel types

- Example
 - multithreaded server creates new agent to handle each incoming request



Process & channel types

```
type result_value is . . . ;
type result_channel is channel result_value;
type result_ptr is access result_channel;

type request_info is record
    . . . ;
    result_please : result_ptr;
end record request_info;

type client is process body
    port ( channel request : out request_info );
    variable result : result_ptr := new result_channel;
begin
    send ( . . . , result ) to request;
    receive . . . from result.all;
end process body client;
```

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Process & channel types (cont)

```
type server is process body
    port ( channel request : in request_info );

    type agent is process body
        port ( channel request : in request_info );
        variable info : request_info;

        begin
            receive info from request;
            . . .
            send . . . to info.result_please.all;
            exit;
        end process body agent;

    . . .
```

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Process & channel types (cont)

```
...  
variable info : request_info;  
channel agent_request : request_info;  
type agent_ptr is access agent;  
variable new_agent : agent_ptr;  
  
begin -- server  
  receive info from request;  
  new_agent := new agent  
    port map ( agent_request );  
  send info to agent_request;  
end process body server;
```

Conclusions

- Improve support for system modeling in VHDL
 - more abstract communication
 - process abstraction
 - process and channel types
 - dynamic creation/termination
- Integrate with existing language
 - extend existing mechanism where appropriate
- Enables single-language design flow
 - hardware/software co-design